## Computational Geometry • 3 January 2023

POLYGONS AND TRIANGULATIONS Def: A polygon is the closed region of the plane bounded by a finite collection of line segments forming a closed curve that does not intersect itself. Not Polygons: Polygons: Let P be a polygon. DIAGONAL: line segment connecting two vertices of P, lying in the interior of P not touching boundary of except at endpoints. TRIANGULATION: A decomposition of P into triangles by a maximal set of noncrossing diagonals can't add any more Elittle theorem, used for proving something else LEMMA: Every polygon with more than 3 vertices has a diagonal. proof: Let v be the lowest vertex of polygon P (righmost, if there are multiple lowest vertices. Let a and b be neighboring vertices ,f v. If segment ab is a diagonal, the we are done. Otherwise, the triangle abov contains some vertex of P. Let L be the line parallel to ab through v. line sweep Sweep L upward toward ab, and let x be the argument first vertex that L meets along this sweep. (choose either, if multiple). Thus, segnent XV is a diagonal of P.

Sample Proof by Induction Prove that  $H2+3+\cdots+n = \frac{n(n+1)}{2}$  for all positive integers n. Call this S(n). Want to prove this is true for  $n = 1, 2, 3, 4\cdots$ Base case: S(1) is  $1 = \frac{1(1+2)}{2}$ , which is clearly true. Induction: Let  $k \ge 1$  be an integer, and assume S(k) is true. [ve want to show that S(k+1) is true.] That is,  $1+2+3+\cdots k = \frac{k(k+1)}{2} \le S(k)$ Add k+1 to both sides:  $1+2+3+\cdots+k+(k+1) = \frac{k(k+1)}{2} = \frac{(k+1)(k+2)}{2}$ want:  $1+2+3+\cdots+k+(k+1) = \frac{(k+1)(k+1+1)}{2} \le S(k+1)$ Thus, by mathematical induction, the statement holds for all positive integers n.

THEOREM: Every polygon has a triangulation. proof: by induction on the number of vertices n of polygon P base case: If n=3, then P is a triangle, which is its own triangulation. induction: Let k=3, and assume that every polygon with fewer than k vertices has a triangulation. [Show that polygons with k vertices have triangulation.] Let P be a polygon with k vertices. By the lemma, P has a diagonal, which splits P into polygons Q and R. Then each of Q and R has fever than k vertices.  $\mathcal{V}$ 

**Theorem:** Every triangulation of a polygon P with n vertices has n-2 triangles and n-3 diagonals.

## Proof (by induction):

First state the base case. Explain why it is true.

If 
$$n=3$$
, then  $P$  is a triangle. A triangulation of a triangle  
has  $n-2=1$  triangles and  
 $n-3=0$  diagonals.

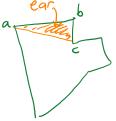
Inductive hypothesis: Let n > 3 be an integer, and assume the statement is true for all polygons with fewer than n vertices.

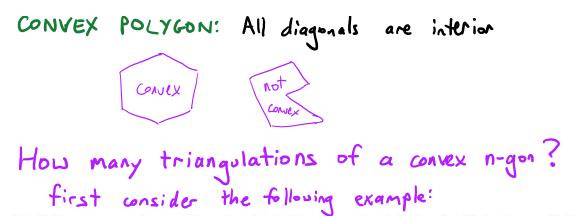
Now explain how the statement follows for a polygon with n vertices.

If P has n vertices, then choose a diagonal (by the lemma)  
that splits P into two polygons: Pi with ni vertices  
and P2 with nz vertices.  
PP, Then 
$$n_1+n_2 = n+2$$
.  
By the induction hypothesis:  
Pi has  $n_1-2$  triangles and  $n_1-3$  diagonals  
P2 has  $n_2-2$  triangles and  $n_2-3$  diagonals.  
Then the number of triangles in P is:  
 $(n_1-2) + (n_2-2) = n_1+n_2-4 = n+2-4 = n-2$ .  
The number of diagonals in P is:  
 $(n_1-3) + (n_2-3) + 1 = n_1+n_2-5 = n+2-5 = n-3$ .  
Note: If a,b,c are consecutive vertices and a

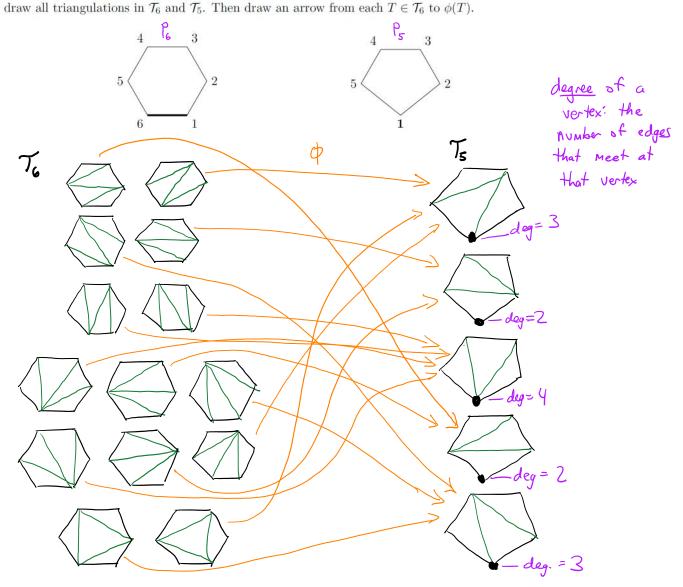
ac is a diagonal, then abc is called an ear.

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Let  $P_n$  be a convex polygon with vertices labeled 1 to n counterclockwise. Let  $\mathcal{T}_n$  be the set of all triangulations of  $P_n$ , and let  $t_n$  be the number of elements of  $\mathcal{T}_n$ . Define the map  $\phi: \mathcal{T}_6 \to \mathcal{T}_5$  that contracts the edge  $\{1, 6\}$  to the point 1. To illustrate this map, first down all triangulations in  $\mathcal{T}$  and  $\mathcal{T}$ . Then down any form each  $T \in \mathcal{T}$  to t(T).



to be continued ...