

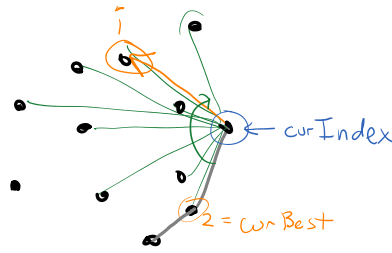
points: stored in a list  
 $[p_1, p_2, \dots, p_n]$

$\text{LeftOf}(a, b, c)$  returns True  
 iff  $c$  is left of  $\vec{ab}$



repeat this  
 for each  
 point in  
 the hull

$O(n)$



### IMPLEMENTATION

Let  $\text{curIndex}$  be index of current point on the hull.

Let  $\text{curBest}$  be the index of some point  
 if  $\text{curIndex} = 1$ , then set  $\text{curBest} = 2$   
 otherwise,  $\text{curBest} = 1$

For  $i$  from 2 to  $n$ :

If  $i = \text{curBest}$  or  $i = \text{curIndex}$ : continue

If  $\text{LeftOf}(\text{pts}[\text{curIndex}], \text{pts}[i], \text{pts}[\text{curBest}])$ :  
 set  $\text{curBest} = i$

Let  $h$  be the number of points in the hull.  
 Then the gift-wrapping algorithm is  $O(nh)$ .

Worst case:  $n=h$ , so alg is  $O(n^2)$ .

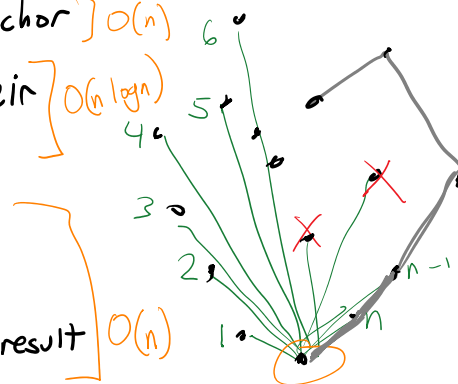
### GRAHAM SCAN ALGORITHM - $O(n \log n)$

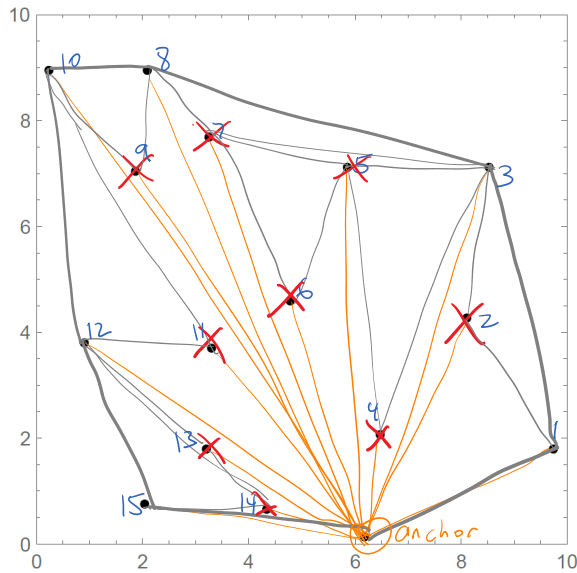
1. Choose the lowest point as an anchor  $O(n)$

2. Sort all the other points by their angle with the anchor  $O(n \log n)$

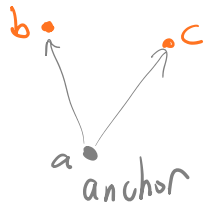
3. Construct the hull as follows:

- append next point to the hull
- remove any reflex vertices that result  $O(n)$





## SORTING POINTS



is  $b$  left of  $\vec{ac}$ ?

If yes, then  $c$  comes before  $b$ .

If no, then  $b$  comes before  $c$ .

Fact: Sorting  $n$  items can be done in  $O(n \log n)$  time, but not faster. ↑ optimal