

# COMPUTATIONAL COMPLEXITY

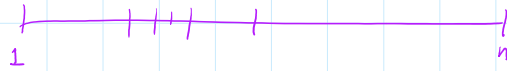
Suppose we have an algorithm that operates on input of size  $n$ . We say the algorithm is  $O(f(n))$  if the runtime is not greater than  $c \cdot f(n)$  for some constant  $c$  (and large  $n$ ).

*"big-O of  $f(n)$ "*

**EXAMPLES:**

$O(1)$  constant time e.g. adding two numbers

$O(\log n)$  log time e.g. binary search — searching for item in a sorted list



$O(n)$  linear time e.g. given list of size  $n$   
 for  $i$  from 1 to  $n$ :  
 Print list[ $i$ ]

$O(n \log n)$  ← fast sorting algorithms (merge sort, quicksort)

$O(n^2)$  ← slow sorting algorithms (bubble sort, insertion sort)

for  $i$  from 1 to  $n$ :  
 for  $j$  from 1 to  $i$ :

## INCREMENTAL ALGORITHM FOR CONVEX HULLS

**INPUT:** set  $S$  of  $n$  points in the plane, given by coordinates

**OUTPUT:** list  $L$  containing vertices of  $\text{conv}(S)$ , in counterclockwise order

**PSEUDOCODE:**

- Sort  $S$  by  $x$ -coordinate
- Take first 3 points. Let  $H_3$  be these points in counterclockwise order

3. For  $k=4$  to  $n$ :

Can we give detailed instructions on how to do this?

Insert point  $k$  into  $H_{k-1}$ , and remove all points now in the interior, forming  $H_k$

4. Set  $L$  equal to  $H_n$

