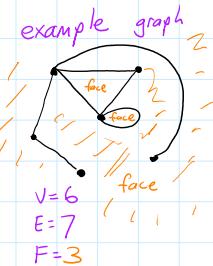


EULER'S FORMULA:

Let G be a connected planar graph with $V \geq 1$ vertices, E edges, and F faces (including the unbounded outer face).

$$\text{Then: } V - E + F = 2$$

proof: Induction on the number of edges.



BASE CASE: If $E=0$, then $V=1$ since graph is connected

$$\text{and } F=1. \text{ So: } V-E+F = 1-0+1 = 2$$

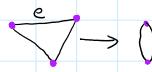
INDUCTION: Assume the formula holds for graphs with up to $E-1$ edges.

Suppose graph G has E edges.

Choose any edge e of G . There are two options:

(1) Edge e connects two different vertices.

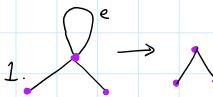
Contract e , reducing V and E by 1.



This does not change the sum $V-E+F$.

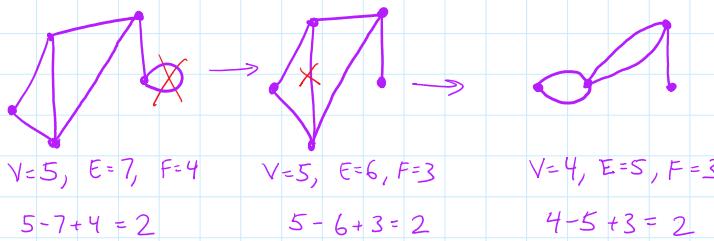
(2) Edge e forms a loop.

Delete edge e , reduces E and F by 1.



Again, this does not change $V-E+F$.

In either case, we reduce G to a graph with $E-1$ edges, without changing the sum $V-E+F$.

EXAMPLE:**TRIANGULATIONS OF POINT SETS**

THEOREM: Let S have h points in its convex hull and k interior points. (Assume not all points are collinear)

Then: any triangulation of S has exactly $2k+h-2$ triangles and $3k+2h-3$ edges.

proof: Let $n=h+k$, the total number of points.

Let t be the number of triangles in a triangulation of S .

$$\text{Then: } V=n$$

$F=t+1$, including the outer face

$$E = \frac{3t+h}{2} \quad \text{numerator counts all edges twice}$$

$$\text{Euler's Formula: } V-E+F = n - \frac{3t+h}{2} + t+1 = 2$$

~ ~ ~ ~ ~ ~ ~ ~ ~ ~ ~ ~ ~

$$2n - 3t - h + 2t + 2 = 4$$

$$2n - t - h = 2$$

solve for t :

$$t = 2n - h - 2$$

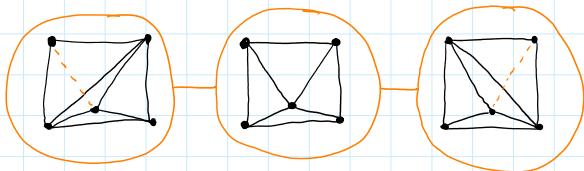
$$t = 2(h+k) - h - 2$$

$$t = 2k + h - 2 \quad \text{number of triangles}$$

number of edges: $E = \frac{3t+h}{2} = \frac{3(2k+h-2)+h}{2} = \frac{6k+4h-6}{2} = \boxed{3k+2h-3}$

FROM ONE TRIANGULATION TO ANOTHER

example:



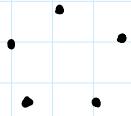
FLIP GRAPH

Vertices: triangulations

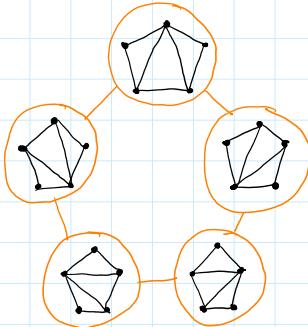
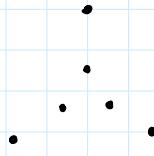
edges: connect two triangulations that differ by a single edge flip

PROBLEMS: Draw the flip graph for each set of points

(a)



(b)



QUESTION: Is the flip graph always connected?