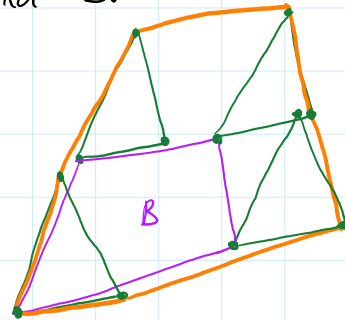
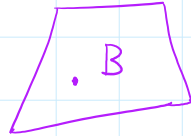
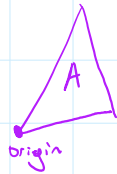


HOW CAN WE COMPUTE THE MINKOWSKI SUM?

Start with convex polygons A and B .

example:



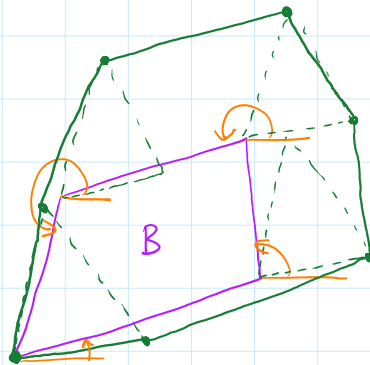
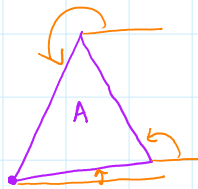
ALGORITHM: For each pair of vertices $v \in A$, $w \in B$, compute the vector sum $v+w$.

Compute the convex hull of all such vector sums.

If A has m vertices and B has n vertices, then we have to consider mn pairs of vertices.

This could be inefficient. Can we do better?

EXAMPLE:



ALGORITHM

Traverse edges of A and B simultaneously, in order of increasing angle with the horizontal axis.

Concatenate traversed edges together to form a polygon, which is $A \oplus B$.

Complexity: $O(m+n)$

WHAT ABOUT NON-CONVEX POLYGONS?

OBSERVE: For any sets A, B, C :

$$A \oplus (B \cup C) = (A \oplus B) \cup (A \oplus C)$$

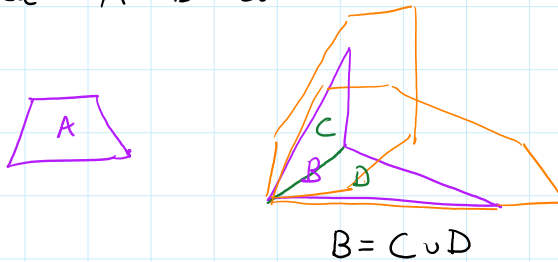
all points in either B or C

vector sum of all $x \in A$ and $y \in B$

vector sum of all $x \in A$ and $y \in C$

vector sum of all $x \in A$ and $y \in (B \cup C)$

Now, suppose A is convex and B is not:

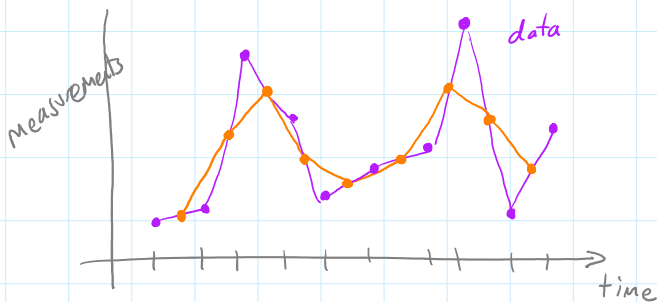


$$A \oplus B = (A \oplus C) \cup (A \oplus D)$$

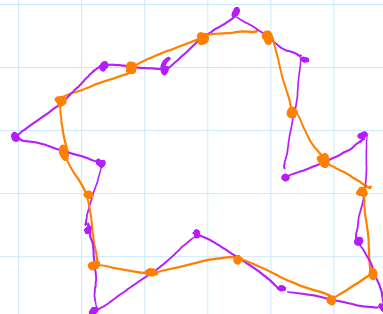
Apply algorithm for convex to each of these, then take union

CURVE SHORTENING

SITUATIONS:



How can we "smooth out" the data?



How can I "smooth out" the polygon?

IDEA: Replace each edge by its midpoint.

MIDPOINT TRANSFORMATION

What happens when we apply the midpoint transformation

... to a polygon?

Same number of vertices.

Length decreases.

... to a polygonal chain?

One less vertex after transformation.

Perimeter is reduced?

Angles get larger — how to make this precise?

If I apply the midpoint transformation to a polygon, do I always get a polygon (without self-intersections)?